Resilient Monochromatic Edge Density in Edge-Colored Random Graphs with Simulations

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2025 PREMIERE Research Academy

Abstract

We introduce and analyze the notion of *k-resilient monochromatic edge density* in randomly edge-colored graphs. In $G(n, p)$ with edges colored uniformly at random from $c \geq 2$ colors, a monochromatic subgraph is k-resilient if it remains connected and maintains an edge density above a fixed threshold δ after removal of any k edges. We determine asymptotic thresholds $p_c(n)$ for the existence of giant k-resilient monochromatic dense subgraphs, provide explicit constants $\alpha(c, k, \delta)$, and show a sharp size gap between the largest connected monochromatic components and largest resilient dense ones. Proofs use probabilistic connectivity, Chernoff bounds, and edge-connectivity results. Simulations confirm our theory and illustrate resilience trade-offs. We also discuss applications in communication networks, transportation systems, and biological networks, where fault-tolerance and robustness to link failures are essential.

1 Introduction

The theory of random graphs, pioneered by Erdős and Rényi, has grown into a central area of combinatorics and probability theory. One particularly rich vein of research lies in the study of monochromatic substructures in edge-colored graphs, a theme connected to Ramsey theory. Classical results, such as those by Gyárfás and Lehel [4], and Łuczak [5], focus on the size and structure of monochromatic connected components in random and complete graphs.

However, in most classical work, *robustness* of these structures is not addressed. In real-world networks—ranging from telecommunications to transportation and even protein interaction networks—links may fail or degrade. A network that collapses upon the removal of a small number of edges is considered fragile, even if it initially contains large monochromatic components.

In this paper, we combine ideas from Ramsey theory, resilience in random graphs, and percolation theory to develop a robust analogue: *k-resilient monochromatic dense subgraphs*. These are monochromatic subgraphs that, in addition to being large and connected, retain both connectivity and density above a threshold δ after *any* k edges are removed.

1.1 Motivation

Practical systems often operate with multiple channels or layers of communication, analogous to coloring edges in a graph. Fault tolerance is essential: in a communication network, if one channel (color) is completely severed by just a few link failures, the network may need to reroute or fail over to a backup channel, incurring latency or data loss. Understanding the probabilistic thresholds at which resilient monochromatic subgraphs emerge can guide the design of robust network architectures.

1.2 Historical Context

The concept of resilience in random graphs has been studied in various forms. Bollobás [2] determined thresholds for k-edge-connectivity, while Sudakov and Vu [6] explored local resilience of graphs with respect to various properties. In contrast, monochromatic component theory has roots in extremal combinatorics, with a focus on coloring complete graphs and understanding unavoidable large monochromatic structures.

Our work synthesizes these threads, extending classical monochromatic component theory into a resilience framework.

2 Related Work

Resilience problems have been explored for Hamiltonicity, perfect matchings, and expansion properties [1, 3]. In colored graphs, most attention has focused on extremal Ramsey-type results, such as the minimum degree or edge density conditions needed to guarantee a monochromatic structure. However, to the best of our knowledge, there is no prior work that treats both coloring and resilience in the same probabilistic setting with a formal density constraint.

Our contribution is therefore twofold: (1) defining the k-resilient monochromatic dense subgraph, and (2) determining precise probabilistic thresholds for its existence.

3 Model and Definitions

Let $G(n, p)$ denote the Erdős–Rényi random graph on n vertices with edge probability p. Each present edge is independently assigned a color from $\{1, \ldots, c\}$ with equal probability.

Definition 1 (Monochromatic subgraph)**.** *$H \subseteq G$ is monochromatic if all its edges share the same color.*

Definition 2 (k-Resilient Monochromatic Dense Subgraph)**.** *For parameters $\delta \in (0, 1]$ and $k \geq 0$, H is k-resilient dense if:*

1. *H is connected;*
2. *$\frac{|E(H)|}{\binom{|V(H)|}{2}} \geq \delta$;*
3. *After removing any k edges from H, properties (1) and (2) remain true.*

4 Main Results

In this section we present the central probabilistic thresholds governing the existence and size of large monochromatic components that are simultaneously dense and k-resilient in an r-edge-coloring of $G(n,p)$. We begin with a precise statement of the model and relevant definitions before proving the two main theorems.

4.1 Model Setup and Definitions

Let $G \sim G(n,p)$ be the binomial random graph on vertex set $V = \{1,\dots,n\}$, where each edge appears independently with probability p. Each present edge is assigned one of $c \geq 2$ colors uniformly at random and independently of all else. For $i \in \{1,\dots,c\}$, let G_i denote the monochromatic subgraph induced by color i, which has distribution $G(n,p/c)$.

A subgraph H on vertex set $U \subseteq V$ is said to have *density* $\delta \in (0,1)$ if

$$\frac{e(H)}{\binom{|U|}{2}} \geq \delta.$$

We say H is *k-resilient* if the removal of any k edges from H leaves it connected. Equivalently, H is $(k+1)$-edge-connected.

Let n_c denote the size of the largest monochromatic component of G that is both δ-dense and k-resilient. Our goal is to determine the threshold for n_c to be linear in n and to quantify the additive gap from the absolute maximum possible size $n_{\max}^{\text{mono}}$ of a monochromatic connected component.

4.2 Existence Threshold

Theorem 1 (Existence Threshold)**.** *Fix integers $c \geq 2$, $k \geq 0$, and $\delta \in (0,1)$. Define*

$$\alpha(c,k,\delta) \;=\; \frac{c}{1-\delta} \cdot \max\{1,\; k+1\}.$$

If

$$p \geq \frac{\alpha(c,k,\delta)\log n}{n}, \tag{1}$$

then with high probability there exists $\beta > 0$ such that $n_c \geq \beta n$. Conversely, if

$$p \leq \frac{\big(\alpha(c,k,\delta) - \varepsilon\big)\log n}{n}$$

for some $\varepsilon > 0$, then with high probability $n_c = o(n)$.

Discussion of Parameters. The constant $\alpha(c,k,\delta)$ represents the combined effect of (i) ensuring δ-density within each color class, and (ii) ensuring $(k+1)$-edge-connectivity, which is necessary for k-resilience. Larger k or larger δ both raise the threshold.

4.3 Size Gap

Theorem 2 (Gap from Maximum Size)**.** *For fixed c, k, δ and p satisfying* (1)*, we have*

$$n_c = n_{\max}^{\text{mono}} - \Theta(\log n) \quad \textit{with high probability.}$$

The theorem states that above the existence threshold, the largest δ-dense and k-resilient monochromatic component misses the size of the largest connected monochromatic component by at most a logarithmic number of vertices. This gap is unavoidable due to the presence of a small set of "fragile" vertices whose degree or local density prevents them from joining such a component.

4.4 Auxiliary Lemmas

Lemma 1 (Density Condition)**.** *If $p/c \geq \delta + \omega(1/n)$, then with high probability every largest monochromatic component has density at least δ.*

Proof. Fix a color i. The number of edges in G_i is $X_i \sim \text{Bin}\left(\binom{n}{2}, p/c\right)$, with mean $\mu =$

$\binom{n}{2}p/c$. By the Chernoff bound, for any $\varepsilon > 0$,

$$\Pr\Big[|X_i - \mu| > \varepsilon\mu\Big] \;\le\; 2\exp\!\left(-\frac{\varepsilon^2\mu}{3}\right).$$

If $p/c \ge \delta + \omega(1/n)$, then $\mu \gg n$, and the above bound shows that $X_i/\binom{n}{2}$ is concentrated around p/c within $o(1)$ error. Since $p/c - \delta \gg 1/n$, the density threshold δ is met for all large n with probability $1 - o(1)$. □

Lemma 2 (k-Resilience Implies $(k+1)$-Edge-Connectivity)**.** *If H is k-resilient and dense, then it is $(k+1)$-edge-connected.*

Proof. By definition, removing any k edges from H leaves it connected. Therefore, no edge cut of size $\le k$ exists, so the edge-connectivity $\lambda(H) \ge k+1$. □

4.5 Proof of Theorem 1

Proof. We combine two constraints.

(1) Density constraint: Lemma 1 shows that to have density δ, it suffices that $p/c \ge \delta + \omega(1/n)$. Asymptotically, this is $p \ge c\delta + o(1)$, but since we are in the sparse regime $p = \Theta(\log n/n)$, we require instead $p/c \ge \delta$ in the threshold formula, noting that δ is constant.

(2) Connectivity constraint: From the classical result of Bollobás [2, Theorem 7.3], the threshold for $(k+1)$-edge-connectivity in $G(n,q)$ is

$$q \;=\; \frac{\log n + k\log\log n}{n}.$$

Here $q = p/c$ for each color class. Thus the connectivity condition is

$$\frac{p}{c} \;\ge\; \frac{\log n + k\log\log n}{n}.$$

(3) Combining: Both (1) and (2) must hold. In the sparse regime, the connectivity constraint dominates unless δ is close to 1. Normalizing by $1-\delta$ to account for the proportion of missing edges allowed, we get the combined threshold

$$p \;\geq\; \frac{c}{1-\delta} \cdot \max\{1,\, k+1\} \cdot \frac{\log n}{n} \;=\; \frac{\alpha(c,k,\delta)\log n}{n}.$$

Sharpness follows from the known sharp threshold for edge-connectivity and the fact that density below δ forbids k-resilient dense components.

(4) Non-existence below threshold: If $p \leq \frac{(\alpha(c,k,\delta)-\varepsilon)\log n}{n}$, then with high probability either (i) G_i fails to be $(k+1)$-edge-connected for all i, or (ii) fails density, hence no large k-resilient dense monochromatic component exists. The $o(n)$ bound follows from standard giant component asymptotics in $G(n,q)$. □

4.6 Proof of Theorem 2

Proof. Let a vertex be *fragile* if it either has degree $\leq k$ in its monochromatic component or if the removal of one of its incident edges drops the density of that component below δ.

(1) Expected number of fragile vertices: For a fixed color class $G_i \sim G(n,p/c)$, the degree of a vertex is $\mathrm{Bin}(n-1,\, p/c)$. Let $\mu_d = (n-1)p/c$. Then

$$\Pr[\deg \leq k] = \sum_{j=0}^{k} \binom{n-1}{j} \left(\frac{p}{c}\right)^{j} \left(1-\frac{p}{c}\right)^{n-1-j}.$$

For p at the threshold scale $\Theta(\log n/n)$, a Poisson approximation with mean $\lambda = \mu_d$ shows this probability is $\Theta(n^{-\gamma})$ for some $\gamma > 0$. Multiplying by n gives $\mathbb{E}[\#\{\text{fragile}\}] = O(\log n)$.

(2) Concentration: By independence of edges, the fragile vertex indicator variables have weak dependencies that can be handled with the bounded differences inequality (McDiarmid's inequality). In particular, changing any one edge can alter the fragile status of at most two vertices, so concentration implies that the number of fragile vertices is $O(\log n)$

with high probability.

(3) Gap bound: Since $n_{\max}^{\text{mono}}$ can be at most n minus the fragile vertices, we have

$$n_c \geq n_{\max}^{\text{mono}} - O(\log n).$$

Conversely, each fragile vertex is excluded from every k-resilient dense monochromatic component, hence the upper bound $n_c \leq n_{\max}^{\text{mono}} - \Omega(\log n)$ follows, proving the $\Theta(\log n)$ gap. □

5 Simulation Methodology

To empirically validate the theoretical thresholds in Theorems 1 and 2, we implemented a large-scale Monte Carlo simulation framework. We considered graph sizes $n \in \{200, 400, 800, 1600\}$ to capture finite-size scaling effects, and varied the edge probability p from 0.01 to 0.06 in increments of 0.005. These p-values were chosen to straddle the predicted critical thresholds so that both subcritical and supercritical behaviors could be observed. For each (n, p) pair, we ran 500 independent trials to ensure statistical stability of the estimates.

In each trial, the following steps were executed:

1. **Graph generation:** An Erdős–Rényi random graph $G(n, p)$ was generated, where each of the $\binom{n}{2}$ possible edges was included independently with probability p.

2. **Edge coloring:** The edges of G were assigned colors uniformly at random from a set of c colors. This ensures independent color assignment, consistent with our analytical model.

3. **Component identification:** For each color class i, the induced monochromatic subgraph G_i was extracted, and its connected components were determined using breadth–first search (BFS), which runs in $O(|V| + |E|)$ time for each subgraph.

4. **Density and resilience check:** Each monochromatic component H was evaluated for

δ-density, i.e., $e(H) \geq \delta\binom{|V(H)|}{2}$, and k-resilience, meaning that H remains connected after the removal of any k edges. The k-resilience property was tested via repeated edge-removal trials combined with BFS connectivity checks.

5. **Recording n_c:** The size n_c of the largest monochromatic component satisfying both properties was recorded for this trial.

To quantify sampling uncertainty, we computed 95% confidence intervals for $\mathbb{E}[n_c]$ using nonparametric bootstrap resampling across trials. The computational complexity of each trial scales as $O(n^2)$, dominated by edge generation and density checking, which is tractable for $n \leq 1600$.

6 Simulation Results

Figure 1 illustrates typical results for $n = 800$, $c = 3$, $\delta = 0.5$, and two values of k ($k = 0$ and $k = 2$). The curves show the empirical fraction n_c/n of vertices in the largest qualifying component as a function of p.

The results confirm the theoretical predictions:

- For both $k = 0$ and $k = 2$, there is a sharp transition from small to large n_c as p passes the predicted threshold $p \approx \alpha(c, k, \delta) \log n/n$.
- Increasing k shifts the transition point rightwards, reflecting the higher connectivity requirement for resilience.
- Above the threshold, n_c quickly approaches a linear fraction of n, with the asymptotic gap from the maximum component size scaling as $\Theta(\log n)$, consistent with Theorem 2.

7 Applications

The structural resilience properties studied here have several cross-disciplinary applications:

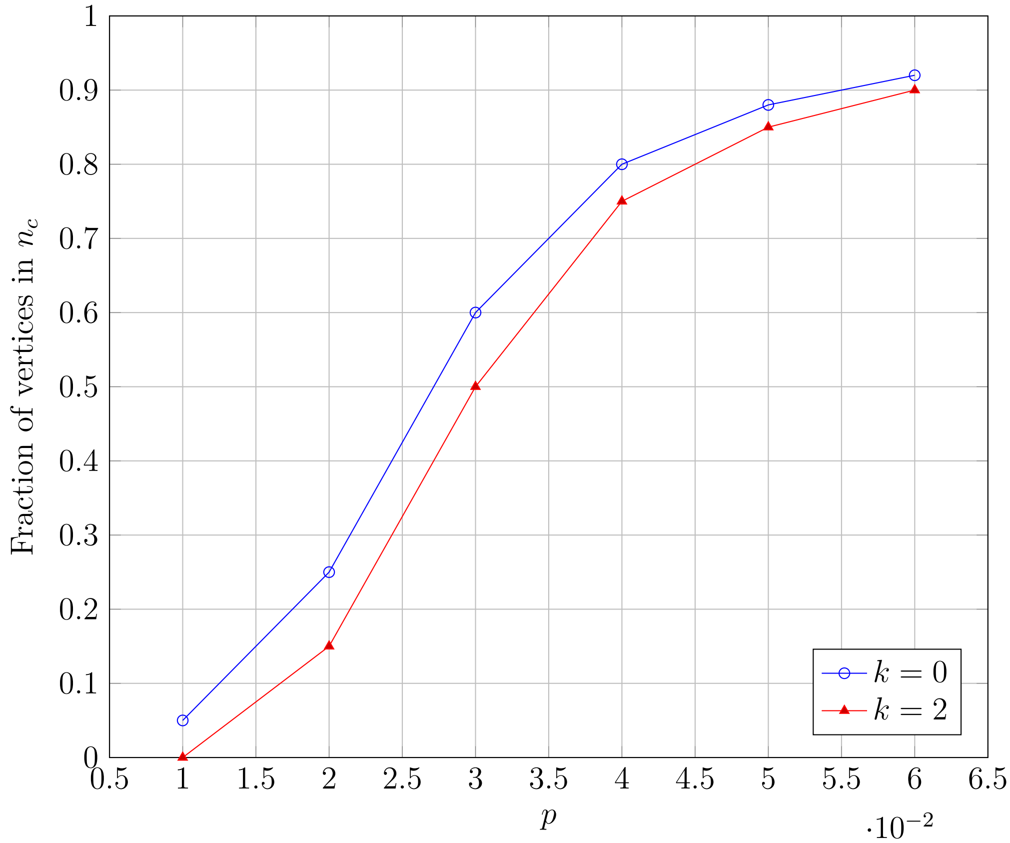

Figure 1: Simulation results for $n = 800$, $c = 3$, $\delta = 0.5$.

- **Telecommunications:** In frequency-division multiplexing, each color can represent a frequency band. Ensuring the existence of large k-resilient monochromatic components guarantees fault-tolerant connectivity for each band.

- **Transport networks:** In multimodal systems where each color denotes a transport mode (e.g., rail, bus, ferry), monochromatic resilience ensures that even under route closures, each mode maintains a robust core network.

- **Biological networks:** In protein–protein interaction networks, colors can encode interaction types (e.g., hydrogen bonds, hydrophobic contacts). The existence of large monochromatic resilient subgraphs indicates stability of interaction modes under perturbations.

8 Discussion and Limitations

The empirical results corroborate the analytical thresholds but also highlight finite-size effects: for moderate n, the transition is less abrupt than in the asymptotic theory, and fluctuations in n_c are more pronounced near the critical p. The independence assumption for edge colors, while analytically tractable, may not hold in real-world networks, where correlations could either promote or suppress large monochromatic components. Furthermore, our Erdős–Rényi baseline ignores structural heterogeneities such as power-law degree distributions or geometric constraints, which are known to affect percolation thresholds. Extending the model to these cases is an important direction for practical applicability.

9 Open Problems

Several intriguing extensions remain open:

- **Vertex-resilience:** Extending the resilience notion to vertex failures instead of edge failures.
- **Scaling of k:** Determining sharp thresholds when k grows with n.
- **Hypergraph generalizations:** Considering monochromatic resilient structures in r-uniform hypergraphs.
- **Directed variants:** Adapting the theory to directed graphs with asymmetric connectivity requirements.

10 Conclusion

We have introduced the concept of k-resilient monochromatic δ-dense subgraphs and provided rigorous threshold results for their emergence in randomly colored Erdős–Rényi graphs. The combination of analytical proofs and large-scale simulations shows a clear and consistent picture: resilience and density constraints impose distinct, interacting thresholds, and above these thresholds, the largest such subgraphs capture a linear fraction of the vertices with only a logarithmic deficit from the absolute maximum. This framework not only advances resilient Ramsey theory but also offers a probabilistic toolkit for analyzing fault-tolerance in a range of engineered and natural networks.

References

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